Randomness extraction from Bell violation with continuous parametric down conversion

Lijiong Shen Jianwei Lee Alessandro Cerè Le Phuc Thinh Valerio Scarani Christian Kurtsiefer



Jean-Daniel Bancal University of Basel, CH

Antia Lamas-Linares Texas Advanced Computing Center, USA

Thomas Gerrits Adriana E. Lita Sae Woo Nam *NIST, USA* 

# Secure private communication requires Randomness

Classical systems cannot guarantee unpredictability



# Secure private communication requires Randomness

Classical systems cannot guarantee unpredictability





# Non-classical correlations certify genuine randomness



Correlation for measurement settings  $\alpha_x$ ,  $\beta_y$ 

$$E_{x,y} = P(a = b | x, y) - P(a \neq b | x, y)$$

Bell parameter from 4 settings

$$S = E_{00} + E_{01} + E_{10} - E_{11}$$



# Our point to the state of the art



S. Pironio et al., Nature (2010)

B. G. Christensen et al., PRL (2013)

Y.Liu et al., PRL (2018)

P.Bierhorst et al., Nature(2018)

# Bell test with CW source and two detectors



# Bell test with CW source and two detectors



# Correlation depends on bin width $\boldsymbol{\tau}$



### Correlation depends on bin width $\tau$

![](_page_9_Figure_1.jpeg)

### Correlation depends on bin width $\tau$

![](_page_10_Figure_1.jpeg)

# Experimental setup

![](_page_11_Figure_1.jpeg)

#### P. H. Eberhard, Phys. Rev.A(1993)

#### Pair generation rate $\approx$ 24 000/s for 5 mW of UV pump

![](_page_12_Figure_1.jpeg)

#### Optimal setting for loophole free Bell test - State

 $\ket{\psi}pprox$  0.9  $\ket{HV}$  - 0.43  $\ket{VH}$ 

![](_page_13_Figure_2.jpeg)

#### Optimal setting for loophole free Bell test - State

 $\ket{\psi}pprox$  0.9  $\ket{HV}$  - 0.43  $\ket{VH}$ 

![](_page_14_Figure_2.jpeg)

### Optimal setting for loophole free Bell test - Projections

![](_page_15_Figure_1.jpeg)

### We observe a violation of S = 2.01602(32)

![](_page_16_Figure_1.jpeg)

# Rate of randomness

![](_page_17_Figure_1.jpeg)

# Rate of randomness

![](_page_18_Figure_1.jpeg)

# Finite statistics - Block extraction

![](_page_19_Picture_1.jpeg)

$$m = n \cdot \eta_{\text{opt}}(\varepsilon_c, \varepsilon_s) + 4 \log rac{\epsilon_{EX}}{n} - 10$$
  
We choose $arepsilon_c = \varepsilon_s = 10^{-10}$ 

Trevisan extractor based on polynomial hashing with block weak design

# In 26 min we generated 617 920 random bits (396 bits/s)

![](_page_20_Picture_1.jpeg)

![](_page_20_Picture_2.jpeg)

# Processing time for 26mins data

٠

26 mins data  $\downarrow$  9 hours processing

# Processing time for 10 hours data

![](_page_22_Figure_1.jpeg)

10 hours data  $\downarrow$  28 years processing

Toeplitz extractor  $\downarrow$  943 bits/s in few hours

# Our point to the state of the art

![](_page_23_Figure_1.jpeg)

S. Pironio et al., Nature (2010)

B. G. Christensen et al., PRL (2013)

Y.Liu et al., PRL (2018)

P.Bierhorst et al., Nature(2018)

# Conclusion

![](_page_24_Figure_1.jpeg)

![](_page_25_Picture_0.jpeg)

# SUZHOU

![](_page_26_Picture_1.jpeg)

# Observed violation changes with $\boldsymbol{\tau}$

![](_page_28_Figure_1.jpeg)

### Rate of randomness with finite block size

![](_page_29_Figure_1.jpeg)